



1. Use the Akkra-Bazi Theorem to derive closed form expressions for the following recursively defined functions:
 - a) $T(n) = 2 \cdot T(n/4) + \sqrt{n}$
 - b) $T(n) = 3 \cdot T(n/2) + n \log n$
 - c) $T(n) = T(n/2) + \alpha$
 - d) $T(n) = 4 \cdot T(n/2) + n^2 / \log n$

It is assumed that in all cases we have $T(n) = c \cdot n$ for $n \leq 4$, and α and c are some positive constants.

2. Use the Akkra-Bazi Theorem to derive the “Master Theorem” for divide-and-conquer recurrence relations, which we stated in class and which you find in most textbooks on analysis of algorithms, stated in some form, e.g. as

Theorem (Master Theorem)
 Let $T(n)$ be a monotonically increasing function that satisfies

$$\begin{aligned} T(n) &= aT\left(\frac{n}{b}\right) + f(n) \\ T(1) &= c \end{aligned}$$

where $a \geq 1, b \geq 2, c > 0$. If $f(n) \in \Theta(n^d)$ where $d \geq 0$, then

$$T(n) = \begin{cases} \Theta(n^d) & \text{if } a < b^d \\ \Theta(n^d \log n) & \text{if } a = b^d \\ \Theta(n^{\log_b a}) & \text{if } a > b^d \end{cases}$$

3. Prove the other inequality of the AB-Theorem stated in class.

Let

$$f(x) = \begin{cases} h(x) & \text{for } 1 \leq x \leq x_0 \\ af(x/b) + g(x) & \text{for } x > x_0. \end{cases}$$

where

- a) $a > 0, b > 1$, and $x_0 \geq b$ are constants,
- b) $x \geq 1$ is a real number,
- c) $d_1 \leq h(x) \leq d_2$ for some positive constants d_1, d_2 and all x with $1 \leq x \leq x_0$, and
- d) g is a nonnegative function satisfying the polynomial growth condition, i.e. there are positive constants c_1 and c_2 such that

$$c_1 g(x) \leq g(u) \leq c_2 g(x) \quad \text{for all } x > x_0 \text{ and } u \in [x/b, x],$$

- e) and the following technical condition holds:

$$I := \int_1^{x_0} \frac{g(u)}{u^{p+1}} du < \infty.$$

Let p be the unique real number for which $a/b^p = 1$. Then

$$f(x) = \Omega\left(x^p \left(1 + \int_1^x \frac{g(u)}{u^{p+1}} du\right)\right).$$